Sample questions for midterm 2
CS 421, Spring 2009

1. Recall that in lectures 11 and 12 we gave several translation schemes for expressions and statements: [e] produces a pair containing the code for e and the location of the result; [S] gives the code for S; $[e]_x$ gives the code to evaluate e and put its value in x; $[e]_{Lt,Lf}$ translates a Boolean expression e to code that branches to either Lt or Lf (this is called the "short-circuit evaluation scheme"), and $[S]_L$ translates S in a context in which L is the label for a break statement.

a. Generate code for the repeat-until statement: "repeat S until e" executes S and tests e, and repeats until e becomes true. Thus, it is equivalent to "S; while !e do S". Do this in two ways: (i) Using the regular scheme [e] to evaluate the condition; and (ii) Using the short-circuit evaluation scheme for e.

b. Generate code for a multiple assignment statement: (x1, x2) = (e1, e2), which does both assignments "in parallel." Note that this is not that same as doing one assignment followed by the other, because variables x1 and x2 may appear in expressions e1 and e2. Use evaluation scheme $[e]_x$ where appropriate.

c. Give two schemes for conditional expression e1 ? e2 : e3, which gives the value of e2 if e1 is true, or e3 if e1 is false. The two schemes you should provide are (a) the standard [e1 ? e2 : e3], and (b) the assignment scheme $[e1 ? e2 : e3]_x$. You should use the short-circuit evaluation scheme for e1.

2. (a) Name the two parts of a compiler's front end.

(b) Name the ~~two~~ three parts of a compiler's back end.

1. AST → IR
2. Optimization
3. Code generation

(c) What are the two outputs of the front end?

3. Name the items in an activation record. = stack frame

4. Give two advantages of the copying garbage collection algorithm over the non-copying (mark-and-sweep) algorithm.
- Compacts - reduces fragmentation, make it easier to allocate large chunks, improves memory perf.
- Touches only reachable data - faster

5. Give two advantages of the non-copying (mark-and-sweep) garbage collection algorithm over the copying algorithm.
- Less memory overhead
- Compacting may be impossible
- Better if there are large objects.

6. Reference counting is not a popular algorithm. What drawback of this algorithm is the reason?
- Cycles in memory

7. In APL, define multmat n which gives an nxn matrix where position i,j has the value i*j.

multmat 4;;
1 2 3 4

Translation schemes -

$[e]$: instruction to calculate e, gives location of result

$[S]$: code for statement S

$[e]_x$: code to calculate e and put value in x

$[S]_L$: code for S in a while or switch stmt where L is label just after the enclosing stmt

$[e]_{Lt,Lf}$: "short-circuit eval" for boolean exps. Code to jump to Lt or Lf, depending on value of e.

[repeat S until e]

= let L1, L2 = newlabel()
  let (il, loc) = [e]
  in
  L1: [S]
      il
      CJUMP loc, L2, L1
  L2:

C: while (e) S

C: do S while(e)

$[\![\text{repeat } S \text{ until } e]\!]_L$ (using short-circuit eval)

$=$ let $L1, L2 = \text{newlabel}()$

~~let (it, loc) = [e]~~

in

$L1:\ [\![S]\!]_{L2}$

$[\![e]\!]_{L2,L1}$

$L2:$

$[e1 \; ? \; e2 : e3]$

$=$ let $L1, L2, L3 =$ newlabel$()$
$(id2, x2) = [e2] \quad (id3, x3) = [e3]$
$x =$ newlocation$()$

$$\left(\begin{array}{l} [e1]_{L1,L2} \\ L1: \; id2 \\ \quad x = x2 \\ \quad \text{JUMP } L3 \\ L2: \; id3 \\ \quad x = x3 \\ L3: \end{array} \right), \; x$$

$[e1 \; ? \; e2 : e3]_x$

$=$ let $L1, L2, L3 =$ newlabels()
in

$[e1]_{L1,L2}$

$L1:$ $[e2]_x$

JUMP $L3$

$L2:$ $[e3]_x$

$L3:$

2 4 6 8
3 6 9 12
4 8 12 16

8. Define the following OCaml functions. *[Exam: we will provide definitions of fold_right and fold_left.]:*

(a) repeat_until: ('a -> bool) -> ('a -> 'a) -> 'a -> 'a. where repeat_until p f x = x, if p x, or f x if p (f x), or f (f x) if p (f (f x)), etc.

(b) sift: ('a -> bool) -> 'a list -> 'a list * 'a list. sift p lis splits lis into a pair of lists (lis1, lis2), with lis1 containing those elements of lis that satisfy p and lis2 the others.

(c) Write sift using fold_right. Specifically, define sift_base and sift_rec so that
fold_right (sift_rec p) lis sift_base = sift p lis

(d) Write an OCaml function that reverses a list, using *fold_right* instead of explicit recursion.

(e) Write a function f such that map f lis returns a list that contains the absolute values of the elements in lis, in the same order. Do not use any library functions in the definition of f.

(f) Using fold_right and no explicit recursion, define a function that concatenates the elements of a string list.

(g) compose_all [f1;f2;…] a = f1 (f2 (… (fn a)…)). Define compose_all *and* say what its type is.

(h) graph_fun f [x1; x2; ...; xn] = [(x1, f x1); (x2, f x2); ...]. Define graph_fun and say what its type is.
`graph_fun`: $(\alpha \rightarrow \beta) \rightarrow \alpha$ list $\rightarrow (\alpha * \beta)$ list, where

9. What does this OCaml program evaluate to:

```
let x = 4
let y = 6
let f y = x+y
let x = 8
in f(y+x)
```

f y = 4 + y ←

= f 14 = 18

10. Suppose the following Java interface is defined:

```
interface FunObj {
   int apply (int) ;
}
```

A Java class might define a function like this:

```
void map (FunObj f, int[] a) {
```

compose_all [f1; f2; ...; fn] a = f1 (f2 (... (fn a) ...))

compose_all : ('a -> 'a) list -> 'a -> 'a

```
let rec
  compose_all fl a =
    match fl with
      [] -> a
    | f :: fl' -> f (compose_all fl' a)
```

compose_all [f1; f2; ...; fn] a = f1 (f2 (... (fn a) ...))

compose_all : ('a -> 'a) list -> 'a -> 'a

- fold_right f [x1; ...; xn] a = f x1 (f x2 (... f (xn a) ...))

compose_all fl a =
= fold_right (fun f r -> f r) fl a

app

compose_all [incr; decr] 0

→ app incr (app decr 0) → decr 0
= -1

(app decr 0) → (fun f r -> f r)

(fun f r -> f r) incr -1 = incr -1 = 0

void map (FunObj f, int[] a){

```
    for (int i=0; i<a.length; i++)
        a[i] = f.apply(a[i]);
}
```

a. Define classes **decrobj** and **sqrobj** that implement **FunObj** so that **map (new decrobj(), a)** decrements each element of **a**, and **map (new sqrobj(), a, n)** squares each element of **a**.

b. Define a class **compose** that implements **FunObj:**

```
class compose implements FunObj {
    FunObj f, g;
    compose ( FunObj f, FunObj g ) {
        this.f = f; this.g = g;
    }

    int apply (int i) {
        return f.apply (g.apply (i));


    }
}
```

that composes function objects, so that, for example, **map(new compose(new sqrobj(), new decrobj()), a, n)** changes every element a[i] to $(a[i]-1)^2$.

```
class decrobj implements FunObj {
    int apply (int n) {
        return n-1;
    }
```

OCaml: let compose f g = fun n -> f (g n)

compose (fun x -> x+1) (fun x -> x*x)

= fun n -> (fun x -> x+1) ((fun x -> x*x) n)